Objectives
1. To gain experience programming with recursively defined data types in ML.
2. To demonstrate an understanding of static semantics by implementing a type checker.
3. To demonstrate an understanding of dynamic semantics by implementing an interpreter.
4. To gain experience setting up deductive systems.
5. To formalize static and dynamic semantics for a new programming language.

Due Date: Monday, March 19, 2018 (at 5pm).
Late submission: You may submit any part (or both parts) of this assignment late (i.e., between 5pm on 3/19 and 5pm on 3/21) with a 15% penalty on the whole assignment.

Machine Details: Complete this assignment by yourself, the programming portion on the following CSEE network computers: c4lab01, c4lab02, ..., c4lab20. These machines are physically located in ENB 220. Do not use any server machines like grad, babbage, sunblast, etc. You can connect to the C4 machines from home using SSH. (Example: Host name: c4lab01.csee.usf.edu Login ID and Password: <your NetID username and password>). You are responsible for ensuring that your programs compile and execute properly on these machines.

Assignment Description
(1) Programming portion: First, correct any problems with your implementation of the sub function from Assignment III. Then, in a directory containing a copy of as3.sml, begin a new file called as4.sml with the command use "as3.sml"; and then implement the following values.

(a) tc : expr -> typ option
This function takes a STERLING expression e and returns NONE iff e is an ill-typed program and SOME t iff e is a well-typed program having type t.

(b) exception stuck
All you need to do for this step is to declare the exception “stuck”. Please read Section 5.2 of the Elements of ML Programming textbook for details on using exceptions in ML.

(c) eval : expr -> expr
This function takes a STERLING expression e and evaluates e for as many steps as possible. If evaluation of e converges to a value v, then eval(e) returns v; if e diverges then so does eval(e). Function eval must raise exception stuck at any point that evaluation gets “stuck” without a value being produced (but note that because STERLING is type safe, only ill-typed expressions can get stuck before becoming values). Evaluation in STERLING is left-to-right and call-by-value.

Throughout this assignment, you may assume that all variable names in expressions being type checked and evaluated are unique, so you never have to alpha-convert expressions.

The file at http://www.cse.usf.edu/~ligatti/pl-18/as4/moreExprs.sml defines a few STERLING expressions, which may help you test your functions.

Sample Executions
- use "as4.sml";
- ...
- use "moreExprs.sml";
- ...
- tc e;
val it = SOME Int : typ option
- tc eBad;
val it = NONE : typ option
- tc e2;
val it = SOME (Arrow (Int,Arrow (Int,Int))) : typ option
- tc e3;
val it = SOME Int : typ option
- eval e;
[infinite loop here, escaped by pressing Control-c]
Interrupt
- eval eBad;

uncaught exception stuck
raised at: as4.sml:46.31-46.36
- eval e2;
val it =
  FunExpr
  (SOME "f",SOME "x",Int,Arrow (Int,Int),
   FunExpr (NONE,NONE,Int,Int,VarExpr "x")) : expr
- eval e3;
val it = IntExpr 1 : expr

Grading and Submission Notes
For full credit, your implementation must (i) be commented and formatted appropriately (as on previous assignments), (ii) use ML features like pattern matching when appropriate, (iii) not define any extra top-level values, (iv) compile and run on the C4 machines with no errors or warnings (except for stuck exceptions raised at appropriate times), (v) not use any library functions, (vi) not use any ML features that cause side effects to occur (e.g., I/O or references/pointers), (vii) not be significantly more complicated than is necessary, and (viii) be reasonably efficient.

The submission process is the same as for Assignment III, except here you’ll submit as4.sml in Canvas. Please remember to include the pledge as an initial comment; not doing so will lower your grade 50%.

(2) Theory portion:
Define static and dynamic semantics for the language L from Assignment III. Assume that all variable names in all expressions under consideration have been made unique through alpha-conversion; hence, you never have to consider contexts containing more than one entry for the same variable. Also, assume that capture-avoiding substitution ([e/x]e’) is already defined for L, so you can just use that notation ([e/x]e’) without defining it.

As always, avoid making the definitions significantly more complicated than necessary. If you get stuck at any point, please explain in prose whatever you’re having trouble formalizing.

Theory-portion Submission Reminders
- Write the following pledge at the end of your submission: “I pledge my Honor that I have not cheated, and will not cheat, on this assignment.” Sign your name after the pledge. Not including this pledge will lower your grade 50%.
- For full credit, turn in a hardcopy (handwritten or printed) version of your solutions.
- Late submissions may be emailed or submitted in hardcopy.
- All emailed submissions, even if sent before the deadline, will be graded as if they were submitted late, i.e., with a 15% penalty.
- If you think there’s a chance you’ll be absent or late for class on the date this assignment is due, you’re welcome to submit solutions early by giving them to me or the TA before or after class, or during any of our office hours.